Competing Bandits in Non-Stationary Matching Markets

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Joint work with

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> RL Workshop, IISc Bangalore Feb, 2024

Outline

Matching Markets—a Multi Agent Multi-Armed Bandit formulation

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 Competition—Collision—Resolution of conflicts

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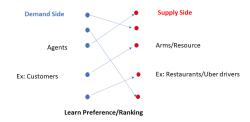
- Matching Markets—a Multi Agent Multi-Armed Bandit formulation
- Competition—Collision—Resolution of conflicts
- Dynamic (Non-Stationary) Markets
 - Algorithm for Non-Stationary Matching Markets
 - Insights for 2 agents
 - & Analysis: Theoretical and Empirical

Competing Bandits/RL in matching markets



Goal: Find Optimal bipartite matching in the presence of competition





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Received a lot of interest in recent years (Liu et.al, 20, Johari et.al, 21, Sankararaman et. al, 21, Basu et.al, 21)





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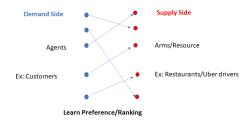
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Motivated by applications in Labor Market (ex. TaskRabbit, Upwork)

- College Admissions (classical motivation, Gale and Shapley, 1962)
- Matching medical interns/residents Dai and Jordan, 2021
- Objective: Match Demand to the Supply side

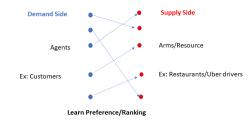
Matching Markets contd.

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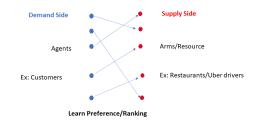


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- Each Agent has a preference over Arms
- Each Arm has preference over Agents

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If these preferences were known, a simple solution is Gale Shapley stable matching algorithm.

However, we don't know these preferences—learn them via successive interactions between Agents and Arms

- Preference modeling via Bandits framework
- We consider N agents and K arms
 - When Agent i pulls arm j-receives a random reward with mean $\mu_{i,j}$
 - Agent i's preference—ordering of the arm means $\{\mu_{i,1}, \mu_{i,2}, \dots, \mu_{i,k}\}$
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Caveat: There is competition in the system—2 or more agents may go for one arm—need to resolve collision/conflicts

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◄ Preferences of Agents (left) are known to the arms (right)-through some common knowledge—one sided learning

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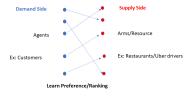
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- & Collision/Competition More than one agents pull same arm–Collision
- Reward is given to the Agent with highest rank among the competitors
- All other Agents receive 0 reward

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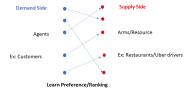
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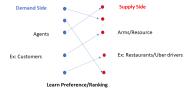
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Under Serial Dictatorship, the matching between Agents and Arms is unique—provides a reference point to characterize Regret

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♣ Let L^(j) be the arm played by an algorithm A.
♣ Regret of agent j playing A is

$$R_j = \sum_{t=1}^{T} \mathbb{E} \left[\mu_{j, \ell_*^{(j)}} - \mu_{j, L^{(j)}} (\mathbf{1}_{L^{(j)} \text{matched player } j}) \right]$$

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Active area of research in Management sciences and Operations Research revolve around understanding the equilibrium properties in such evolving market (Lam et. al 05, Akbarpour et. al 20, Kurino 20, Johari et. al 21)

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We take a step towards bridging the two aforementioned lines of work

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We use smooth varying framework of Wei and Srivastaba '18, Krishnamurthy and Gopalan '21)

 $\clubsuit |\mu_{.,.,t+1} - \mu_{.,.,t}| \leqslant \delta$ for all agent arm pairs

🐥 Maximum drift is δ

Useful in cloud computing, financial applications-decisions in seconds, distribution changes slowly

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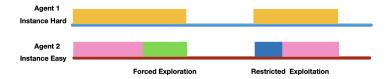
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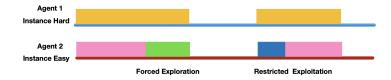
Forced Exploration and Restrictive Exploitation



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Agent 2 will face additional regret. We characterize this regret.

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 \blacklozenge Let us now look at the simplest non-trivial problem with N = 2 agents

Non-stationary Competing Bandits (NSCB) for Agent 1

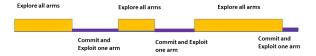
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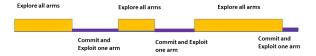
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Since arm means are changing, best arm is changing, so Agent 1 Explore and Commits successively over time.

No market aspect, no competition

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Exploit period such that the arm remains optimal

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- After getting the optimal arm, Agent 1 updates the black-board
- Agent 1 writes the optimal arm index and the exploit time

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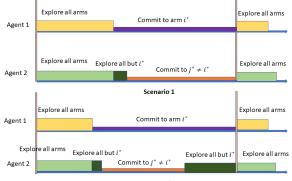
- This is quite common is Distributed Optimization (Wainwright et. al)
- Mode of communication across various agents
- & In Sankararaman et.al '21, this is achieved through structured collision
- In Kong et. al '23, equivalent model is assumed through broadcasting
- Since we do not allow collision, we let the agents use the blackboard to communicate
- Note that this is still a decentralized system, at each time, we allow one agent to update
- After getting the optimal arm, Agent 1 updates the black-board
- Agent 1 writes the optimal arm index and the exploit time
- We later remove the blackboard

- Borns out the competitive nature of the market
- Agent 2-avoids collision; otherwise get 0 reward
- Let us look at different scenarios of Agent 2

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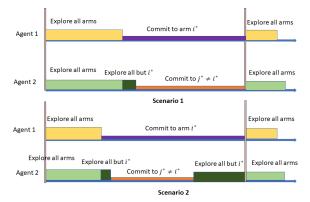


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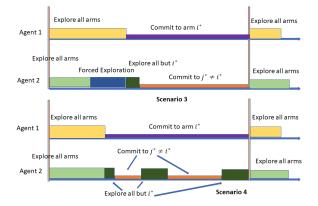


(Domination 1) Restricted Exploration

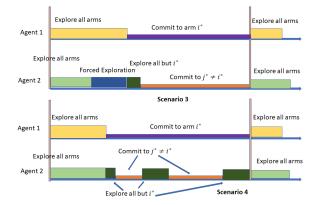
(Domination 2) Commitment to a dominated set of arms

Let us look at different scenarios of Agent 2

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(Domination 3) Forced Exploration

4 (Domination 4) Restricted Exploitation

Let us summarize the actions of Agent 2

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Algorithm 2 NSCB with N = 2; for Agent 2 1: Initialize set of tuples $S_2^{(j)} = \phi, \forall j \in [k]$, episode index $i_2 \leftarrow 1$ 2: for $t = 1, 2, \ldots, T$ do Pull-Arm by Agent 2: 3: if Agent 1 is not committed then 4: Play round robbin on [k] (pull arm t + 1 % k), set 5. $z_t(2) \leftarrow \text{Explore ALL}$ else if Agent 1 is committed to arm j and $S_2^{(j)} = \phi$ 6: then Play round robbin on $[k] \setminus \{j\}$ (i.e., pull arm index 7: t%(k-1)-th smallest arm id in $[k] \setminus \{i\}$ $z_t(2) \leftarrow \text{Explore-j}$ 8: else if Agent 1 is committed to arm j, and $\exists (x, s) \in$ Q٠ $S_2^{(j)}$ s.t. s > t then Play arm x; set $z_t(2) \leftarrow \text{Exploit}$ 10: end if 11: if $\{z_t(1) \neq z_{t-1}(1)\} \in \{z_t(2) \neq z_{t-1}(2)\}$ then 12: 13. $i_2 \leftarrow i_2 + 1, t_{i_2} \leftarrow t$ 14: end if 15: Test by Agent 2: $z_t(2)$ for F [k]s.t. ∈ 16: {Explore-j,Explore ALL} do if Arm $a = \text{Lambda-Opt}(\tilde{\lambda}, [k] \setminus \{j\})$ then $\tau_{ia}^{(j)}$ \leftarrow t - t_{i_2} , buffer₂ 18: $\max\left(\frac{8}{\delta}\sqrt{(k-1)\log T/\tau_{i_2}^{(j)}} - 2(k-2), 0\right)$ if buffer₂ > $\tau_{i_2}^{(j)}$ then 19: $S_2^{(j)} \leftarrow S_2^{(j)^2} \cup \{(a, \min\{t_{i_2} + \mathsf{buffer}_2, s_{i_1+1})\}$ 20: 21: end if 22: end if 23: end for

Problem Complexity–Dynamic Gap

& Dynamic gap for agent r determines how complex the problem is

Average gap between the pairwise arm-means over a window

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Average gap between the pairwise arm-means over a window

Definition

For ${\mathfrak C}\subseteq [k],$ the dynamic gap of Agent r on a dominated set ${\mathfrak C}$ as,

$$\lambda_{t}^{\mathcal{C}}[r] = \max_{\lambda \in [0,1]} \left\{ \min_{\substack{a,b \in [k] \setminus \mathcal{C} \\ a \neq b}} \frac{1}{w(\lambda)} | \sum_{t'=s}^{t} \mu_{r,a,t'} - \mu_{r,b,t'} | \geqslant \lambda \right\},$$

and if such a λ does not exist, we set $\lambda^{\mathcal{C}}_t[r] = c_1 \sqrt{\frac{\log T}{t}}.$ Here, $s = t - w(\lambda) + 1$, and $w(\lambda) = \frac{c_0(k - |\mathcal{C}|)\log T}{\lambda^2}.$

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No superscript when C = φ
 When δ = 0, reduced to the usual gap in MAB
 Generalization of the standard gap

Regret Guarantee

& For Agent 1, regret of successive ETC

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Theorem: Suppose we run NSCB with 2 Agents upto horizon T. Then the expected regret for Agent 1 is $R_1 \lesssim \sum_{l=1}^m \frac{1}{\lambda_{min,l}[1]} k \log T$ and for Agent 2 is $R_2 \lesssim \sum_{l=1}^m \left[\left(\frac{1}{\lambda_{min,l}[2]} + \frac{1}{(\lambda_{min,l}[1])^2} + \frac{1}{min_\alpha} \lambda_{min,l}^{\{\alpha\}}[2] \right) k \log T \right]$

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 $\begin{array}{l} \clubsuit \ T \ is \ divided \ into \ m \ blocks \ with \ length \ \delta^{-2/3} k^{1/3} \log^{1/3} T \\ \clubsuit \ \lambda_{min,l}[r] = min_{t \in l-th \ block} \ \lambda_t[r], \ \lambda_{min,l}^{\{\alpha\}}[r] = min_{t \in l-th \ block} \ \lambda_t^{\{\alpha\}}[r] \end{array}$

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♣ Regret matches that of UCB-D3 in stationary setup by putting $\delta = 0$. We get a regret of $O(\frac{1}{Gap^2} k \log T)$

$\ensuremath{\texttt{NSCB}}$ for N Agents

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When Agent r commits to an arm, it updates the blackboard by the arm index and the committed period

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Algorithm 3 NSCB for r-th Agent

```
1: Input: Horizon T, drift limit \delta
 2: Initialize S_r^{(\Omega)} = \phi for all \Omega \subset [k], and |\Omega| < r - 1,
     Initialize i_r \leftarrow 1, C_1(r) = \phi
 3: RANK ESTIMATION()
 4: for t = 1, 2, ..., T do
 5:
       Update State z_i(r):
       if |\mathcal{C}_t(r)| < r-1 then
 7:
           z_t(r) \leftarrow \text{Explore} - C_t(r)
        else if \exists (x,s) \in S_r^{(\mathcal{C}_t(r))} s.t. s > t, then
 9
 9:
           z_t(r) \leftarrow \text{Exploit}(x)
10:
       else
           z_t(r) \leftarrow \text{Explore} - C_t(r)
11:
12: end if
13: if z_t(r) \neq z_{t-1}(r) then
       i_r \leftarrow i_r + 1, t_{i_-} \leftarrow t
14.
15:
    end if
16: Pull-Arm by Agent r:
17:
     if z_t = \text{Explore} - C_t(r) then
           Play round robbin with [k] \setminus C_t(r) (i.e., pull t + (r - r)
18:
           |\mathcal{C}_t(r)| - 1)\%(k - |\mathcal{C}_t(r)|) smallest arm in [k] \setminus \mathcal{C}_t(r)
19:
       else if z_t = \text{Exploit}(x) then
20:
           Play arm x
       end if
21:
22:
       Test by Agent r:
23:
       for \Omega \subseteq [k] s.t. |\Omega| = r - 1 and z_t(r) = \text{Explore}
       C_t(r) do
           if Arm a = \text{Lambda-Opt}(\tilde{\lambda}, [k] \setminus \Omega) then
24:
              \tau_i^{(\Omega)}
25-
                                        t - t_{i_r}, buffer
              \max\left(\frac{8}{\delta}\sqrt{(k-|\mathcal{C}_t(r)|)\frac{\log T}{r^{(1)}}}-2(k-r),0\right),
              define \overline{t}_{i_r} = \min\{t_{i_r} + \text{buffer}_{i_r}, t_{i_{r-1}+1}\}
                 if buffer<sub>i</sub> > \tau_{i_{-}}, then S_{r}^{\Omega} \leftarrow S_{r}^{\overline{\Omega}} \cup \{(a, \overline{t}_{i_{-}})\},\
26:
              else i_{i_r} \leftarrow i_r + 1, t_{i_r} \leftarrow t
27:
                 end if
           end if
28
29:
       end for
30:
        Updates Black Board:
31: if \exists (x,s) s.t. s \ge t+1, then write (x, \bar{t}_i, r) on the
        black board end if
32:
       Updates Dominated set C_{t+1}(r):
33: Updates C_{t+1}(r) = \{x \in [k] : \exists s > t+1, \text{ and } \exists j < t \in [k]\}
        r-1 s.t. (x, s, j) exists on board}
34: end if
```

```
25. and for
```

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Theorem: Suppose we run NSCB with N Agents upto horizon T. Then the expected regret for Agent r is $R_r \lesssim \sum_{\text{phases}} \left[\left(\frac{1}{\text{Gap}[r]} + \frac{1}{\text{Gap}^2[r-1]} \right) \ k \log T \right]$

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Agent r is dominated by Agents 1,..., r – 1–captured in gap def.
Regret of Agent r depends on all the agents 1,..., r – 1
Similar to the 2 Agent case, the first term comes from Exploring all arms, and the second term comes from forced exploration and restricted exploitation.

With blackboard, Agent 2 knows whether Agent 1 is exploring or committed to a particular arm

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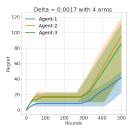
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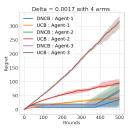
Same idea can be extended to the N Agent setup

Simulations

A Consider 3 Agents N = 3, 4 Arms K = 4



& Compare with Dominated UCB of Sankararaman et. al'21



Dominated UCB is for static market-hence large regret

Open Problems

A Markets and Bandits framework has several open problems

- **♣** We only consider smooth changes with known δ , what about abrupt change or total budgeted change?
- Beyond Serial Dictatorship?

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Spurious Markets: Handling adversaries amidst competition?

Structured Markets: Linear/Contextual Bandits?

Reference

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Thank You