# Competing Bandits in Non-Stationary Matching Markets 

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Joint work with

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## Outline

■ Matching Markets—a Multi Agent Multi-Armed Bandit formulation

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■ Competition-Collision-Resolution of conflicts

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- Matching Markets—a Multi Agent Multi-Armed Bandit formulation

■ Competition-Collision-Resolution of conflicts

- Dynamic (Non-Stationary) Markets
\& Algorithm for Non-Stationary Matching Markets
\& Insights for 2 agents
\% Analysis: Theoretical and Empirical

Matching Markets

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## Competing Bandits/RL in matching markets



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\& Motivated by applications in Labor Market (ex. TaskRabbit, Upwork)
\& College Admissions (classical motivation, Gale and Shapley, 1962)
\& Matching medical interns/residents Dai and Jordan, 2021
\& Objective: Match Demand to the Supply side

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\% If these preferences were known, a simple solution is Gale Shapley stable matching algorithm.
\& However, we don't know these preferences-learn them via successive interactions between Agents and Arms

Model via Bandits framework

## Model via Bandits framework

\& Preference modeling via Bandits framework
\& We consider N agents and K arms

- When Agent $i$ pulls arm $j$-receives a random reward with mean $\mu_{i, j}$
- Agent $i$ 's preference—ordering of the arm means $\left\{\mu_{i, 1}, \mu_{i, 2}, \ldots, \mu_{i, k}\right\}$
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2. Solving multi-agent multi armed Bandit problem is equivalent to learning preferences
\& Use tools from bandit literature-Liu et.al, 20, Jagadeesan et. al, 21, Johari et.al, 21, Sankararaman et. al, 21, Basu et.al, 21

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\& Use tools from bandit literature-Liu et.al, 20, Jagadeesan et. al, 21, Johari et.al, 21, Sankararaman et. al, 21, Basu et.al, 21
\& Caveat: There is competition in the system-2 or more agents may go for one arm—need to resolve collision/conflicts


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\& Collision/Competition More than one agents pull same arm-Collision
\& Reward is given to the Agent with highest rank among the competitors
\& All other Agents receive 0 reward

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\& Under Serial Dictatorship, the matching between Agents and Arms is unique-provides a reference point to characterize Regret

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\& Let $L^{(j)}$ be the arm played by an algorithm $\mathbb{A}$.
\& Regret of agent $j$ playing $\mathbb{A}$ is

$$
\mathrm{R}_{\mathrm{j}}=\sum_{\mathrm{t}=1}^{\mathrm{T}} \mathbb{E}\left[\mu_{\mathrm{j}, \ell_{*}^{(j)}}-\mu_{\mathrm{j}, \mathrm{~L}(\mathrm{j})}\left(\mathbf{1}_{\mathrm{L}(\mathrm{j})} \text { matched player } \mathfrak{j}\right)\right]
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\& With this, regret of $\mathcal{O}(j \mathrm{~K} \log \mathrm{~T})$ is obtained for $j$-th ranked agent


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\& We take a step towards bridging the two aforementioned lines of work

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\% We use smooth varying framework of Wei and Srivastaba '18, Krishnamurthy and Gopalan '21)
\& $\left|\mu_{\ldots, \ldots+1}-\mu_{\ldots, ., t}\right| \leqslant \delta$ for all agent arm pairs
\& Maximum drift is $\delta$
\& Useful in cloud computing, financial applications-decisions in seconds, distribution changes slowly

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$\measuredangle$ Agent 2 will face additional regret. We characterize this regret.

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© Let us now look at the simplest non-trivial problem with $\mathrm{N}=2$ agents

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\& Since arm means are changing, best arm is changing, so Agent 1 Explore and Commits successively over time.
« No market aspect, no competition

## NSCB for Agent 1 contd.

\& Epoch based $-s_{1}, s_{2}, .$. denote starting of epoch
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$\triangleleft$ Exploit period such that the arm remains optimal


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\% Mode of communication across various agents
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- We later remove the blackboard


## NSCB for Agent 2

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\& Agent 2-avoids collision; otherwise get 0 reward
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\& (Domination 3) Forced Exploration
\& (Domination 4) Restricted Exploitation

## NSCB for Agent 2 contd.

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```
Algorithm 2 NSCB with \(N=2\); for Agent 2
    Initialize set of tuples \(S_{2}^{(j)}=\phi, \forall j \in[k]\), episode index
    \(i_{2} \leftarrow 1\)
    for \(t=1,2, \ldots, T\) do
        Pull-Arm by Agent 2:
        if Agent 1 is not committed then
            Play round robbin on \([k]\) (pull arm \(t+1 \% k\) ), set
            \(z_{t}(2) \leftarrow\) Explore ALL
        else if Agent 1 is committed to arm \(j\) and \(S_{2}^{(j)}=\phi\)
        then
            Play round robbin on \([k] \backslash\{j\}\) (i.e., pull arm index
            \(t \%(k-1)\)-th smallest arm id in \([k] \backslash\{j\})\)
            \(z_{t}(2) \leftarrow\) Explore-j
        else if Agent 1 is committed to arm \(j\), and \(\exists(x, s) \in\)
        \(S_{2}^{(j)}\) s.t. \(s>t\) then
            Play arm \(x\); set \(z_{t}(2) \leftarrow\) Exploit
        end if
        if \(\left\{z_{t}(1) \neq z_{t-1}(1)\right\} \bigcirc \mathcal{R}\left\{z_{t}(2) \neq z_{t-1}(2)\right\}\) then
            \(i_{2} \leftarrow i_{2}+1, t_{i_{2}} \leftarrow t\)
        end if
        Test by Agent 2:
        for \(j \in[k]\) s.t. \(\quad z_{t}(2) \quad \in\)
        \{Explore-j, Explore ALL\} do
            if Arm \(a=\) Lambda-Opt \((\tilde{\lambda},[k] \backslash\{j\})\) then
                    \(\tau_{i_{2}}^{(j)} \leftarrow t-t_{i_{2}}, \quad\) buffer \(_{2} \quad=\)
            \(\max \left(\frac{8}{\delta} \sqrt{(k-1) \log T / \tau_{i_{2}}^{(j)}}-2(k-2), 0\right)\)
            if buffer \({ }_{2}>\tau_{i_{2}}^{(j)}\) then
                \(S_{2}^{(j)} \leftarrow S_{2}^{\left(j^{2}\right.} \cup\left\{\left(a, \min \left\{t_{i_{2}}+\right.\right.\right.\) buffer \(\left.\left._{2}, s_{i_{1}+1}\right)\right\}\)
                end if
            end if
        end for
```


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For $\mathcal{C} \subseteq[k]$, the dynamic gap of Agent $r$ on a dominated set $\mathcal{C}$ as,

$$
\lambda_{t}^{\mathcal{e}}[r]=\max _{\lambda \in[0,1]}\left\{\min _{\substack{a, b \in[k] \backslash e \\ a \neq b}} \frac{1}{w(\lambda)}\left|\sum_{t^{\prime}=s}^{t} \mu_{r, a, t^{\prime}}-\mu_{r, b, t^{\prime}}\right| \geqslant \lambda\right\}
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and if such a $\lambda$ does not exist, we set $\lambda_{t}^{e}[r]=c_{1} \sqrt{\frac{\log T}{t}}$. Here, $s=t-w(\lambda)+1$, and $w(\lambda)=\frac{c_{0}(k-|\mathcal{C}|) \log T}{\lambda^{2}}$.

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$s=t-w(\lambda)+1$, and $w(\lambda)=\frac{c_{0}(k-|\mathcal{C}|) \log T}{\lambda^{2}}$.
\& No superscript when $\mathcal{C}=\phi$
\& When $\delta=0$, reduced to the usual gap in MAB
\& Generalization of the standard gap

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\mathrm{R}_{1} \lesssim \sum_{\mathrm{l}=1}^{\mathrm{m}} \frac{1}{\lambda_{\min , l}[1]} \mathrm{k} \log \mathrm{~T}
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and for Agent 2 is
$R_{2} \lesssim \sum_{l=1}^{m}\left[\left(\frac{1}{\lambda_{\text {min }, l}[2]}+\frac{1}{\left(\lambda_{\text {min }, l}[1]\right)^{2}}+\frac{1}{\min _{\mathrm{a}} \lambda_{\text {min }, \mathrm{l}}^{\{\mathrm{a}\}}[2]}\right) \mathrm{k} \log \mathrm{T}\right]$

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$$

\& $T$ is divided into $m$ blocks with length $\delta^{-2 / 3} \mathrm{k}^{1 / 3} \log ^{1 / 3} T$
\& $\lambda_{\text {min, }}[r]=\min _{t \in l-\text { th block }} \lambda_{t}[r], \lambda_{\text {min }, l}^{\{a\}}[r]=\min _{t \in l-\text { th block }} \lambda_{t}^{\{a\}}[r]$

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- Third Term: Comes from Restrictive exploration (Explore-j) phase, where Agent 1 is committed to arm $j$
\& Regret matches that of UCB-D3 in stationary setup by putting $\delta=0$.
We get a regret of $\mathcal{O}\left(\frac{1}{\text { Gap }^{2}} k \log T\right)$

NSCB for N Agents

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\& Also, Agent $r$ ends faces restricted exploitation if anyone of agents $\{1, \ldots, r-1\}$ ends their exploitation.
\& When Agent r commits to an arm, it updates the blackboard by the arm index and the committed period

## NSCB for N Agents

\% Let us summarize the actions of Agent $r$

## NSCB for N Agents

## \& Let us summarize the actions of Agent r

```
Algorithm 3 NSCB for r}r\mathrm{ -th Agent
    1: Input: Horizon T, drift limit }
    2: Initialize }\mp@subsup{S}{r}{(\Omega)}=\phi\mathrm{ for all }\Omega\subseteq[k]\mathrm{ , and }|\Omega|\leqr-1
    Initialize }\mp@subsup{i}{r}{}\leftarrow1,\mp@subsup{\mathcal{C}}{1}{}(r)=
    3: RANK ESTIMATION()
    for }t=1,2,\ldots,T\mathrm{ do
        Update State }\mp@subsup{z}{t}{}(r)
        if }|\mp@subsup{\mathcal{C}}{t}{}(r)|<r-1 the
        z
        else if }\exists(x,s)\in\mp@subsup{S}{r}{(\mp@subsup{\mathcal{C}}{t}{\prime}(r))}\mathrm{ s.t. }s>t\mathrm{ , then
            zt(r)\leftarrowEExploit}(x
        else
            zt}(r)\leftarrow\mathrm{ Explore - }\mp@subsup{\mathcal{C}}{t}{}(r
    end if
    if }\mp@subsup{z}{t}{}(r)\not=\mp@subsup{z}{t-1}{}(r)\mathrm{ then
        irr}\leftarrow\mp@subsup{i}{r}{}+1,\quad\mp@subsup{t}{\mp@subsup{i}{r}{}}{}\leftarrow
    end if
    Pull-Arm by Agent r:
    if }\mp@subsup{z}{t}{}=\mathrm{ Explore - - }\mp@subsup{\mathcal{C}}{t}{}(r)\mathrm{ then
            Play round robbin with [k]\\mp@subsup{\mathcal{C}}{t}{}(r) (i.e., pull }t+(r
            |\mathcal{C}
        else if }\mp@subsup{z}{t}{}=\operatorname{Exploit}(x)\mathrm{ then
            Play arm x
        end if
        Test by Agent r:
        for \Omega\subseteq[k] s.t. }|\Omega|=r-1 and \mp@subsup{z}{t}{}(r)=\mathrm{ Explore-
        \mathcal{C}
            if Arm a = Lambda-Opt ( }\tilde{\lambda},[k]\\Omega)\mathrm{ then
            \mp@subsup{\tau}{\mp@subsup{i}{r}{}}{(\Omega)}\leftarrowt<t-t, buffer,
            max}(\frac{8}{\delta}\sqrt{}{(k-|\mp@subsup{\mathcal{C}}{t}{}(r)|)\frac{\operatorname{log}T}{\mp@subsup{\tau}{ir}{(\Omega)}}}-2(k-r),0)
            define }\mp@subsup{\overline{t}}{\mp@subsup{i}{r}{}}{}=\operatorname{min}{\mp@subsup{t}{\mp@subsup{i}{r}{}}{}+\mp@subsup{\mathrm{ buffer r}}{\mp@subsup{i}{r}{}}{},\mp@subsup{t}{\mp@subsup{i}{\Gamma}{}-1+1}{}
                if buffer }\mp@subsup{i}{\mp@subsup{r}{r}{}}{}>\mp@subsup{\tau}{\mp@subsup{i}{r}{}}{\prime}\mathrm{ , then }\mp@subsup{S}{r}{\Omega}\leftarrow\mp@subsup{S}{r}{\Omega}\cup{(a,\mp@subsup{\overline{t}}{\mp@subsup{i}{r}{}}{})}\mathrm{ ,
            else }\mp@subsup{i}{\mp@subsup{i}{r}{}}{}\leftarrow\mp@subsup{i}{r}{}+1,\mp@subsup{t}{\mp@subsup{i}{r}{}}{}\leftarrow
                end if
        end if
    end for
    Updates Black Board:
if }\exists(x,s) s.t. s\geqt+1, then write (x,\mp@subsup{t}{\mp@subsup{i}{r}{}}{},r)\mathrm{ on the
    black board end if
    Updates Dominated set }\mp@subsup{\mathcal{C}}{t+1}{}(r)\mathrm{ :
33: Updates }\mp@subsup{\mathcal{C}}{t+1}{}(r)={x\in[k]:\existss>t+1, and \existsj
    r-1 s.t. (x,s,j) exists on board}
    end if
```

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\& Agent $r$ is dominated by Agents $1, \ldots, r-1$-captured in gap def.
\& Regret of Agent $r$ depends on all the agents $1, \ldots, r-1$
\& Similar to the 2 Agent case, the first term comes from Exploring all arms, and the second term comes from forced exploration and restricted exploitation.

## Learning without Blackboard; $\mathrm{N}=2$

\& With blackboard, Agent 2 knows whether Agent 1 is exploring or committed to a particular arm
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\& Toggling $Q_{t}$ is sufficient to get the information broadcasted by the blackboard
\& Same idea can be extended to the N Agent setup


## Simulations

\& Consider 3 Agents $\mathrm{N}=3,4$ Arms $\mathrm{K}=4$

\& Compare with Dominated UCB of Sankararaman et. al'21

\& Dominated UCB is for static market-hence large regret

## Open Problems

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\& Spurious Markets: Handling adversaries amidst competition?
\& Structured Markets: Linear/Contextual Bandits?

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## Thank You

